## Some Selected Picks from Extremal Combinatorics

Sam Spiro, UC San Diego.

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For example, a central problem of Ramsey Theory is to determine how large a clique $K_{N}$ one can have before every two-coloring of its edges contains a monochromatic clique of size $t$. I.e., it seeks to determine how many "vertices" a "two-colored cliques $K_{N}$ without a monochromatic clique of size $t$ " can have.

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- AMGM:

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\left(\prod_{i=1}^{n} x_{i}\right)^{1 / n} \leq n^{-1} \sum_{i=1}^{n} x_{i}=\mathbb{E}\left[x_{i}\right]
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## Extremal Graph Theory

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How many "edges" can an " $n$-vertex $F$-free graph" have?

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## Question

How many "edges" can an " $n$-vertex $F$-free graph" have?
We denote this quantity by ex $(n, F)$, which is called the extremal number or Turán number of $F$.

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For example, ex $\left(n, K_{3}\right) \geq\left\lfloor n^{2} / 4\right\rfloor$ because $G=K_{\lfloor n / 2\rfloor,\lceil n / 2\rceil}$ is an $n$-vertex graph which is $K_{3}$-free and with $\left\lfloor n^{2} / 4\right\rfloor$ edges.

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Theorem (Mantel, 1907)

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\operatorname{ex}\left(n, K_{3}\right)=\left\lfloor n^{2} / 4\right\rfloor
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and the unique construction is $K_{\lfloor n / 2\rfloor,\lceil n / 2\rceil}$.

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## Theorem (Turán, 1941)

$$
\operatorname{ex}\left(n, K_{t}\right)=\left\lfloor\left(1-\frac{1}{t-1}\right) \frac{n^{2}}{2}\right\rfloor
$$

and the unique construction is a complete $(t-1)$-partite graph with parts of sizes $\lfloor n /(t-1)\rfloor,\lceil n /(t-1)\rceil$.

## Extremal Graph Theory

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## Extremal Graph Theory

Turán's Theorem can be generalized considerably. Recall that $\chi(G)$ is the chromatic number of $G$, i.e. the fewest number of colors needed to properly color the vertices of $G$ (meaning that every pair of adjacent vertices receives a distinct color).
For example, the following shows that the Petersen graph $G$ has $\chi(G) \leq 3$, and it is not hard to show that it is in fact equal to 3 .


Picture form Wikipedia.

## Extremal Graph Theory

## Theorem (Erdős-Stone, 1946)

For any graph F,

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\operatorname{ex}(n, F)=\left(1-\frac{1}{\chi(F)-1}\right) \frac{n^{2}}{2}+o\left(n^{2}\right)
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Note that when $\chi(F)=2$, this only gives

$$
\operatorname{ex}(n, F)=o\left(n^{2}\right)
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## Extremal Graph Theory

In general determining ex $(n, F)$ when $F$ is bipartite (i.e. $\chi(F)=2$ ) is very, very difficult, even for simple families of graphs.

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## Theorem

If $F$ is a tree on $k$ vertices, then

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Conjecture (Erdős-Sós, 1962)
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## Extremal Graph Theory

Theorem (Kovari-Sós-Turán, 1954)
If $s \leq t$, then

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\operatorname{ex}\left(n, K_{s, t}\right) \ll n^{2-1 / s} .
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This is known to hold whenever $t>(s-1)$ ! due to the existence of projective norm graphs.

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In general, it is often the case that a conjectured upper bound is relatively easy to obtain, but seemingly the lower bound requires tools from algebra, geometry, number theory, etc.

## Hypergraphs

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When each $e \in E$ has $r$ vertices we say that $H$ is an $r$-uniform hypergraph.

## Hypergraphs

If $F$ is an $r$-uniform hypergrpah, define ex $(n, F)$ to be the maximum number of edges an $n$-vertex $r$-uniform hypergraph $H$ can have without containing $F$ as a subgraph.

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Very little is known about ex $(n, F)$ for hypergraphs. We do not even know the asymptotic value of ex $\left(n, K_{4}^{(3)}\right)$, where $K_{4}^{(3)}$ is the complete 3 -uniform hypergraph on 4 vertices (which is the simplest non-trivial clique that is not a graph).

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In general, many problems which are easy for graphs become much harder for hypergraphs.

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and in general the hypergraph $S_{V}^{(r)}$ is called a star.

## Hypergraphs

## Theorem (Erdős-Ko-Rado, 1961)

If $n \geq 2 r$ then

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\operatorname{ex}\left(n, M_{2}^{(r)}\right)=\binom{n-1}{r-1}
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## Conjecture (Erdős Matching Conjecture, 1965)

If $n \geq s r-1$, then

$$
\operatorname{ex}\left(n, M_{s}^{(r)}\right)=\max \left\{\binom{n}{r}-\binom{n-s+1}{r},\binom{s r-1}{r}\right\} .
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## Combinatorial Number Theory

We say that a set of $k$ integers is a $k$-term arithmetic progression (or simple a $k$-AP) if they are of the form $\{a, a+d, a+2 d, \ldots, a+(k-1) d\}$.

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## Question

How many "integers in $\{1, \ldots, n\}$ " can "a set $A$ not containing any $k-A P "$ have?

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Let us adopt the (non-standard) notation ex( $n, k-\mathrm{AP}$ ) for the size of a largest $k$-AP free subset of $\{1, \ldots, n\}$.

## Combinatorial Number Theory

## Theorem (Behrend, 1946)

There exists a constant $c>0$ so that for $k \geq 3$,

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\operatorname{ex}(n, k-A P) \geq e^{-c \sqrt{\log n}} \cdot n .
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For all $k$,

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Theorem (Szemerédi, 1975)
For all $k$,

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That is, every dense set of integers contains arbitrarily long arithmetic progressions.

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Theorem (Schoen, May 2020)

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## Theorem (Schoen, May 2020)

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$$

This upper bound is roughly $e^{-\log \log n} \cdot n$, while the best lower bound is still $e^{-c \sqrt{\log n}} \cdot n$, and closing this a gap is a huge problem in the field.

## Erdős Numbers

As you may have realized, Erdős has written a lot of papers (both in extremal combinatorics and in general).

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As you may have realized, Erdős has written a lot of papers (both in extremal combinatorics and in general).

This has led to the amusing statistic: Erdős has Erdős number 0. If you write a paper with someone who has Erdős number $k$, then your Erdős number is $k+1$.

## Erdős Numbers

Almost every professional mathematician has a relatively small Erdős number (e.g. in the single digits).

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| MR Erdos Number $=4$ <br> Claus Mazanti Sorensen |
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It tends to be even smaller once you start picking extremal combinatorialists!

MR Erdos Number $=2$

| Sam Spiro | coauthored with | Fan Chung | MR4057523 |
| :--- | :--- | :--- | :--- |
| Fan Chung | coauthored with | Paul Erdős $^{1}$ | MR0889356 |

## Wrapping Things Up

We've discussed a lot of the major results and problems in extremal combinatorics.

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Rapping Things Up

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Bear in mind for this R A P,

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Bear in mind for this R A P, I got a Behrend mind and no K A P.

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Am I getting my point across, should I spell it out for you?

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Am I getting my point across, should I spell it out for you? AMGM, all your products are less than my average, $\mu$.

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Getting 10 out of 10 's Fall, Winter, Spring and Summer, Your only 10 is your Erdős number.

## Rapping Things Up

Now I'd rather be nice, stay home and count,

## Rapping Things Up

Now I'd rather be nice, stay home and count, But if you want to fight, you're going down for the count.

## Rapping Things Up

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Bring your whole clique, I'll bring my acrobatics, Color you black and blue until you're monochromatic.

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Don't think I'm care-free just cause I'm doing mathematics, If you mess with me, things will get hyper graphic.

The End

## (Drops Mic)

## Answers to Questions

1 Yes, I did make the rap first and then designed the talk around explaining it.

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2 Yes, I am open to performing at weddings as well as bar/bat mitzvahs.
3 My (tentative) rapper name is "Lil Oh."

